

AP COMPUTER SCIENCE A – RECURSIVE BINARY SEARCH

The two examples which follow concern the sorted list of 18 numbers shown at lower left. The variables b , e and m are the current positions of the beginning, end, and middle of the part of the list that currently is being searched. Note that, initially, for iteration 1, $b = 0$ and $e = 17$. Also, $m = (b + e)/2$. As we discussed last semester, the binary search algorithm has on order of $\log n$ operations, where n is the length of the list. Here, $\log_2 18 = 4.17$ so that the maximum number of iterations possible is 5.

i	i	i
n 1	t	t
d i	e	e
e s	r Example 1:	r Example 2:
x t	a -----	a -----
--+--	t	t
0 17	i	i
1 23	o	o
2 25	n b m e	n b m e
3 29	---+---+---+---	---+---+---+---
4 33	1 0 8 17	1 0 8 17
5 37	2 8 12 17	2 8 12 17
6 41		3 8 10 12
7 42		4 10 11 12
8 48		5 10 10 11
9 54		
10 59		
11 65		
12 73		
13 78		
14 82		
15 85		
16 90		
17 97		

Example #1. Here we are searching for 73. For iteration 1, we check if $\text{list}[m]$, or $\text{list}[8] = 48$, is equal to 73. It is not, so since $48 < 73$, b is redefined as $m = 8$. For iteration 2, $\text{list}[12] = 73$, so the number 73 is found at index 12 after 2 iterations.

Example #2. Here we are searching for 60. For iteration 1, we check if $\text{list}[8] = 48$ is equal to 60. It is not, so since $48 < 60$, b is redefined as $m = 8$. For iteration 2, we check if $\text{list}[12] = 73$ is equal to 60. It is not, so since $60 < 73$, e is redefined as $m = 12$. For iteration 3, we check if $\text{list}[10] = 59$ is equal to 60. It is not, so since $59 < 60$, b is redefined as $m = 10$. For iteration 4, we check if $\text{list}[11] = 65$ is equal to 60. It is not, so since $60 < 65$, e is redefined as $m = 11$. For iteration 5, we check if $\text{list}[10] = 59$ is equal to 60. It is not. Since $b == m$, the process is stopped (and the number has not been found). In general, the process should be stopped if $(b == m) \ || \ (m == e)$.